Emulation based analysis using binary instrumentation

Application on CTF





SPEAKERS

Myunghun Cha

- From Republic of Korea
- POSTECH senior student majoring CSE
- Team Leader of PLUS
- CODEGATE 2009 Hacking Contest 3rd place
- DEFCON 2009 CTF 3rd place
- DEFCON 2011 CTF 8th place
- Many hacking contest experience

Jinsuk Park

- POSTECH sophomore majoring ME
- Team member of PLUS

PLUS

- POSTECH Laboratory for UNIX Security
- Found in 1992
- Researching on various security issues
- Participating in lots of hacking contests
- Participated in DEFCON CTF three times
 - -2009 (3rd)
 - -2010 (3rd)
 - -2011 (8th)
 - -2012

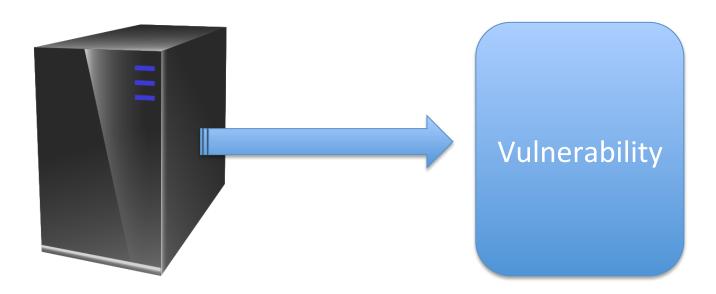


Motivation

DEFCON CTF

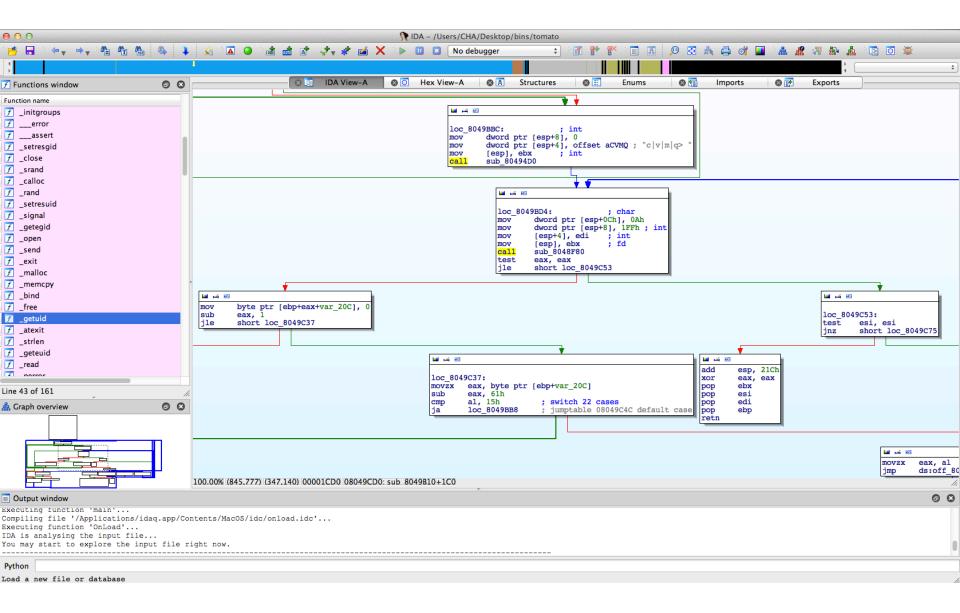
CTF Basic Rule

- CTF: Capture The Flag
- Each team is given vulnerable server
- Vulnerable daemons are running on the server



CTF Daemon

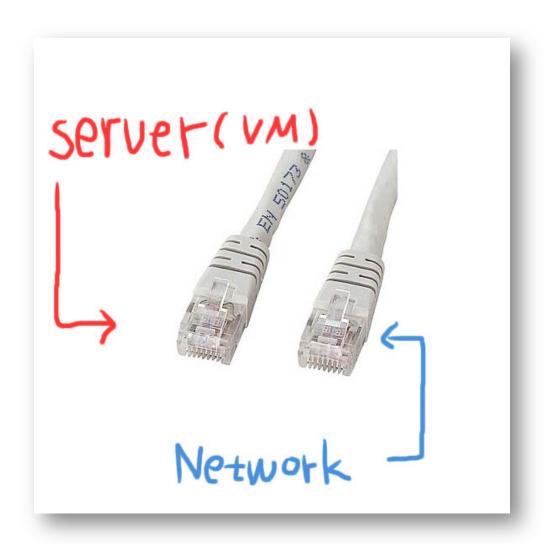
```
plus# nc localhost 6391
Fresh Tomatoes! The real scoop on the movie poop!
MY MALLOC
MY MALLOC
TY MALLOC
MY MALLOC
civia> c
Please to be making the comments!
this is comment
MY MALLUC
Tanks, mi tink!
civimiq> m
Which of these baller babies needs to be replaced?
asdf
Replacing this one:
What uu wanna renlace it wit?
asdfasdf
civimiq>
```

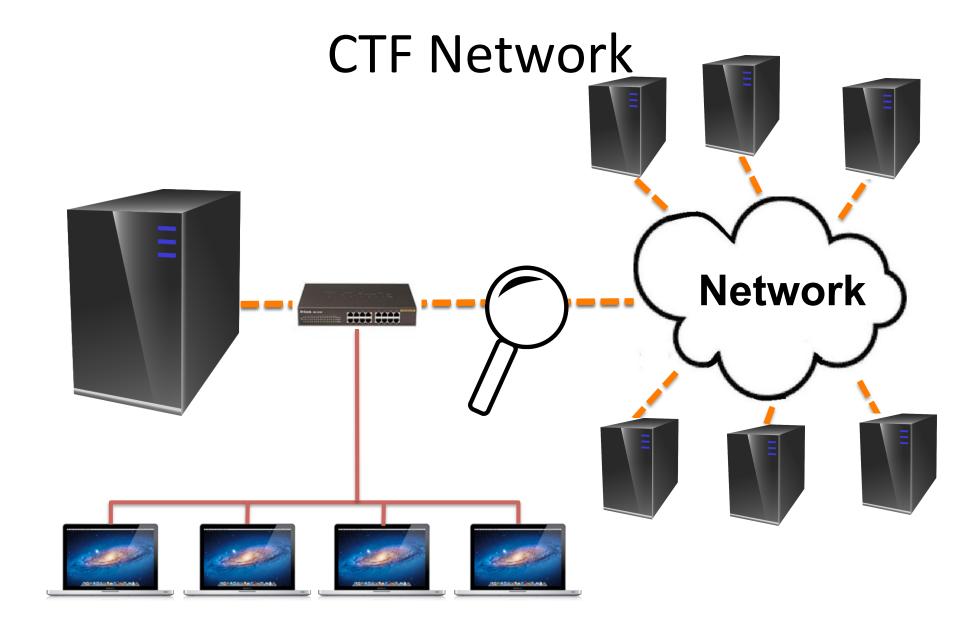


Scoring

- There's a key file for each daemon which is changed periodically
- You should read or write the key file to get a score
- It simulates information stealling and corruption in real world

CTF Network





CTF Summary

We can attack over the wire

We can sniff, suspect, or drop packet

We can attack analyzing binary
 or using other teams' exploit



What do I want to do?

I want to detect attacks

 I want to analyze vulnerability easily using other teams' attack

• Then... how?

EMULATION BASED ANALYSIS

Emulation Based Analysis

- We can detect bug following specific patterns
 - Stack boundary check
 - memcpy without string length check
 - EIP address check
 - Format string from user input
- Verification user input is much more easier than finding hidden bug
- Dynamic analysis is easier than static analysis

Instrumentation?



Dynamic Binary Instrumentation

 Ability to monitor or measure the level of a program's <u>performance</u>, to <u>diagnose errors</u> and to <u>write trace</u> <u>information</u>



Dynamic Binary Instrumentation

 A technique to analyze and modify the behavior of a binary program by injecting arbitrary code at arbitrary places while it is executing

Usage

- Simulation / Emulation
- Performance Analysis
- Program optimization
- Bug detection
- Correctness Checking
- Call graphs
- Memory Analysis

For hackers?

- Fuzzing
- Covert Debugging
- Exploitable Vulnerability Detection
- Automated Reverse Engineering
- Bypass Anti-Debuggers
- Automated exploitation
- Automated unpacking

DBI frameworks

- Pin
- Valgrind
- DynamoRio
- Etc.

Why valgrind?

 Valgrind runs on BSD but PIN does not (which is DEFCON CTF Environment)

Valgrind: Introduction

- Valgrind Core
 - DBI framework
 - Simulated CPU
- Valgrind tool
 - Written in C using Valgrind framework
 - Used as Plug-ins for Valgrind
- Valgrind Core + tool plug-in = Valgrind tool

Valgrind: Tools

- Memcheck: check memory management of the binary executable
- Cachegrind: cache profiling
- Helgrind: Data races conditions detection
- Massif: Heap profiler
- User written tool
- usage: valgrind --tool=<toolname> [options] prog-and-args

Valgrind: How it works

1. Disassembly

Machine code(x86)

↓
Intermediate Language(IR)

2. Instrumentation

↓ Instrumented IR

IR

3. Assembly

Instrumented IR

↓

Machine

code(x86)

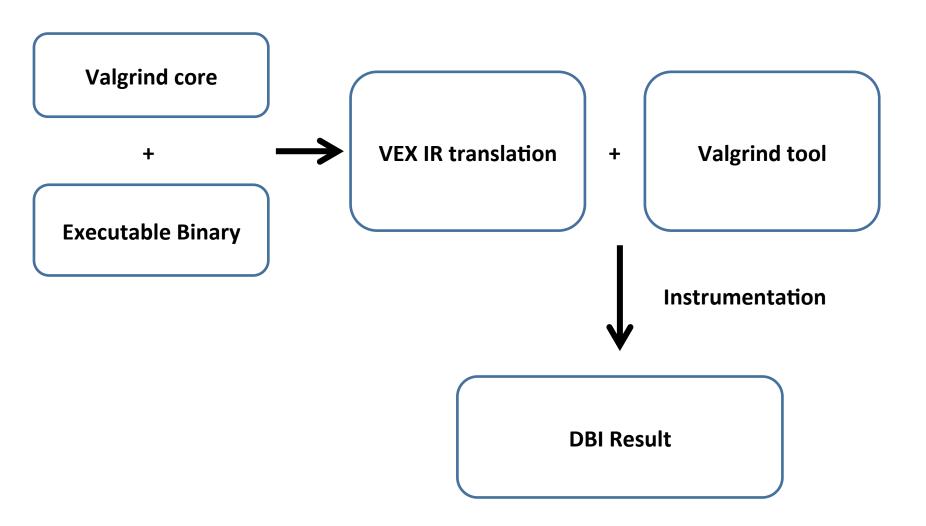
VEX IR(Intermediate Representation)

- Valgrind's binary translation mechanism
- VEX IR: machine independent intermediate representation
- Translates a block of binary codes to simplified VEX representation

VEX IR: Example

- addl %eax, %ebx :
 - -t3 = GET:132(0) # get %eax, a 32-bit integer
 - -t2 = GET:132(12) # get %ebx, a 32-bit integer
 - -t1 = Add32(t3,t2) # addl
 - -PUT(0) = t1 # put %eax

Valgrind: Overview



Attack Detection Using Valgrind DBI Framework

CTFGRIND

What does it do?

- match registered execution patterns
- checks sensitive memory area overwriting
- marks execution flow using IDA Plug-in

Pattern 1: RET overwriting

- We can get the guest machine's register values
- We should protect our RET and stored EBP

- 1. Monitor every memory operation (Store)
- 2. Compare target address with \$EBP
- 3. Output callstack

Pattern 2: GOT overwriting

 We can do in the same manner, because the address of GOT is static in a binary

Pattern 3: Strcpy

 What if a bug comes from using library function such as strcpy

- We can compare the RET before the library function call and after the call
- 2. There could be many vulnerable library functions such as memcpy, strcpy, and scanf

Possible usage #1

- Attach directly to running daemon
- Prevent attack before exploitation
- Stop the process when a danger is detected
- Possible slow down

Possible usage #2

- Runs on a separated shadow machine
- When it detects attack, register the packet pattern to firewall to prevent further attack
- Can't defend the first attack

IDA Plugin

- CTFGRIND logs the call stack when the attack detected
- IDA Plugin reads the file and marks the execution path
- Helpful to analyze other teams' exploit

DEMO

REFERENCE

- Emulationbased Security Testing for Formal Verification (Bl ack Hat Europe 2009) – Bruno Luiz
- Optimizing binary code produced by Valgrind Luis Veiga
- Valgrind Mario Sanchez, Cecilia Gonzalez
- Hacking using Binary instrumentation Gal Diskin
- Valgrind: A Framework for Heavyweight Dynamic Binary Instrumentation -Nicholas Nethercote, Juliam Seward
- Valgrind Technical Manual

감사합니다

hoon0612 @postech.ac.kr jinmel@postech.ac.kr